| PART VII  INTERNET PROTOCOL: FORWARDING IP DATAGRAMS |
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|   | Datagram Transmission  |
|---|--|
| • | Host delivers datagrams to directly connected machines           |
| • | Host sends datagrams that cannot be delivered directly to router |
| • | Routers forward datagrams to other routers                       |
| • | Final router delivers datagram directly                          |
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| Question                                     |  |
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| Does a host need to make forwarding choices? |  |
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| Question                                   |     |
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| Does a host need to make forwarding choice | es? |
| Answer: YES!                               |     |
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## **Example Host That Must Choose How To Forward Datagrams** path to some path to other destinations Note: host is singly homed!

## Two Broad Cases

- Direct delivery
  - Ultimate destination can be reached over one network
  - The "last hop" along a path
  - Also occurs when two communicating hosts both attach to the same physical network
- Indirect delivery
  - Requires intermediary (router)

| unsmission of an IP datagram between two machines on<br>gle physical network does not involve routers. The send<br>capsulates the datagram in a physical frame, binds th |
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| stination IP address to a physical hardware address, and the resulting frame directly to the destination.  |
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## **Testing Whether A Destination Lies On The Same Physical Network** As The Sender Because the Internet addresses of all machines on a single network include a common network prefix and extracting that prefix requires only a few machine instructions, testing whether a machine can be reached directly is extremely efficient.

| • Ge |  |
|------|--|
|      | neral paradigm   |
| _    | Source host sends to first router                            |
| _    | Each router passes datagram to next router                   |
| -    | Last router along path delivers datagram to destination host |
| • On | ly works if routers cooperate                                |
|      |  |
|      |  |

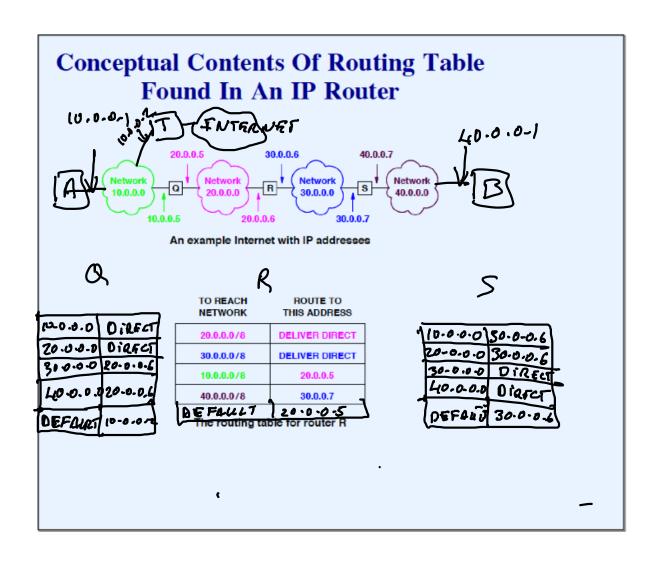
| General Concept |                      |   |
|-----------------|----------------------|---|
| interconnec     | ted structure. Datag | ternet form a cooperative<br>rams pass from router to route<br>n deliver the datagram directly. |
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|                 |                      |   |
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| <ul> <li>Decisions based on table lookup</li> <li>Routing tables keep only network portion of addresses (siz proportional to number of networks, not number of hosts)</li> <li>Extremely efficient         <ul> <li>Lookup</li> <li>Route update</li> </ul> </li> </ul> | oles keep only network portion of addresses (size |
|---|---|
| proportional to number of networks, not number of hosts)  • Extremely efficient  - Lookup   |   |
| – Lookup  | al to number of networks, not number of hosts)    |
|   | efficient   |
| <ul> <li>Route update</li> </ul>  |   |
|   | ıpdate  |
|   |   |
|   |   |
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|   | Important Idea  |
|---|---|
| • | Table used to decide how to send datagram known as routing table (also called a forwarding table) |
| • | Routing table only stores address of next router along the path                                   |
| • | Scheme is known as next-hop forwarding or next-hop routing  |

## **Terminology**

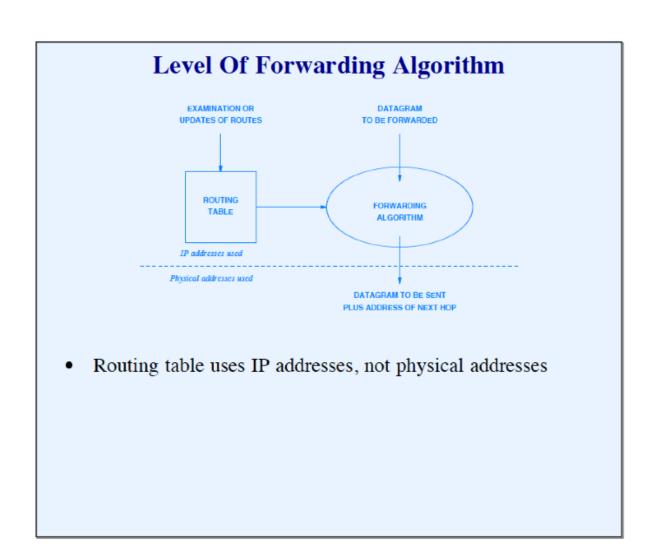
- Originally
  - Routing used to refer to passing datagram from router to router
- More recently
  - Purists decided to use forwarding to refer to the process of looking up a route and sending a datagram
- But...
  - Table is usually called a routing table



| $S_{l}$                                 | pecial Cases |  |
|---|--------------|--|
| Default route                           |              |  |
| <ul> <li>Host-specific route</li> </ul> |              |  |
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|   | Default Route                            |
|---|--|
| • | Special entry in IP routing table        |
| • | Matches "any" destination address        |
| • | Only one default permitted               |
| • | Only selected if no other match in table |
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|             | <b>Host-Specific Route</b>  |
|-------------|---|
| Entry in ro | uting table   |
| Matches en  | ntire 32-bit value  |
|             | ed to send traffic for a specific host along a th (i.e., can differ from the network route) |
| More later  | in the course   |
|             |   |
|             |   |
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|   | Summary   |
|---|---|
| • | IP uses routing table to forward datagrams                      |
| • | Routing table   |
|   | <ul> <li>Stores pairs of network prefix and next hop</li> </ul> |
|   | Can contain host-specific routes and a default route            |
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